# Doubly-linked lists Iterators

Lecture 15 by Marina Barsky

## Doubly-linked Lists: Node

class Node {
 int data;
 Node next;
}



class DoublyLinkedNode {
 int data;
 Node prev;
 Node next;
}

# Doubly-Linked List with tail pointer

- Keeps reference/links in both directions
- Traversing can start from either end

#### DoublyLinkedList:



# In a [doubly-linked] list *head* will be equal to *tail*:

- A. Always
- B. Never
- C. When the list is empty
- D. When there is one element
- E. More than one of the above



### Doubly-Linked List: tradeoffs

- Links in both directions:  $\rightarrow$  can traverse forwards and backwards!
- ALL tail operations (including *remove last*) are fast! Why?
   We have direct access to the tail node & its predecessor
- Additional code complexity in each list operation
   Example: add (int index, E element) need to consider 4 cases: empty list add to front add to front add to tail add in middle
- × Additional space consumption (storing previous)



#### Stitching new node between two existing nodes

The code below adds a new node with data 'X' between two nodes P (parent) and C (child) in a doubly-linked list

```
DoublyLinkedNode x = new DoublyLinkedNode('X');
```

```
if (C != null) C.previous = x;
if (P != null) P.next = x;
```

What should happen if both N and P are null?

- A. Nothing should happen: the code above already covers this case
- B. We need to set head = x;
- C. We need to set C = x;
- D. We need to set P=x;
- E. Something else



Why would anyone use a singly-linked vs. a doubly-linked list?

- A. A singly-linked list uses less memory.
- B. It is easier to implement the insertion at position *i*.
- C. It's faster to remove an element from the end.
- D. None of the above.
- E. More than one of the above



#### Moving heads and tails

- When we add/remove in front we need to update *head*
- When we add/remove at the end we need to update *tail*
- When the linked list currently is or becomes empty: *head=tail*=null
- Many special cases arise!

```
newNode = new DoublyLinkedNode(newData, prev=null, next=null)
if head == null: //empty list
    head = newNode
    tail = head
```





```
newNode = new DoublyLinkedNode(newData, prev=null, next=null)
if head == null: //empty list
    head = newNode
    tail = head
else: //list with at least one real node
    newNode.next = head
    head.prev = newNode
```



```
newNode = new DoublyLinkedNode(newData, prev=null, next=null)
if head == null: //empty list
    head = newNode
    tail = head
else: //list with at least one real node
    newNode.next = head
    head.prev = newNode
    head = newNode
```



#### Sentinel Nodes (aka Dummy nodes)

- We can get rid of special cases if we add fake head and tail nodes
- These are called *sentinel* nodes are they are always present and contain no data
- We can have one sentinel for both head and tail, or we can have a separate node for each
- The head and tail pointers never move and the nodes are inserted between them

Doubly-Linked List with two sentinels: *constructor* 



newNode = new DoublyLinkedNode (newData, prev=null, next=null)

//empty list
newNode.prev = head



newNode = new DoublyLinkedNode (newData, prev=null, next=null)

//empty list
newNode.prev = head
newNode.next = head.next



newNode = new DoublyLinkedNode (newData, prev=null, next=null)

//empty list
newNode.prev = head
newNode.next = head.next
head.next.prev = newNode



newNode = new DoublyLinkedNode (newData, prev=null, next=null)

//empty list
newNode.prev = head
newNode.next = head.next
head.next.prev = newNode
head.next = newNode



newNode = new DoublyLinkedNode (newData, prev=null, next=null)

newNode.prev = head newNode.next = head.next head.next.prev = newNode head.next = newNode

This also works for non-empty lists – there are no special cases!



### Lab 4. Doubly-linked lists with two sentinels

- In Lab 4 you will implement this idea
- Always draw the list before and after each operation to make sure you update all the links correctly



# List Iterators

## Recap ADT *List*: supported operations

ADT *List* supports the following main operations

- Get element by position: get(int index)
- Search element: indexOf(E element)
- Add new element: add(int index, E element)
- Remove element by position: remove(i)

For some problems however these operations are insufficient and we **need access to the underlying implementation of the data** 

### Example: count occurrences

• Write a method *count* that counts the number of times a particular element *o* appears in a List:

```
public static int count(List list, E o) {
    int counter = 0;
    for (int i=0; i<data.size(); i++) {
        E obj = data.get(i);
        if (obj.equals(o)) counter++;
     }
    return counter;
}</pre>
```

 Question: would this work well no matter if the List is an Array List or a Linked List?

### Example: count occurrences

 Write a method *count* that counts the number of times a particular element *o* appears in a List:

```
public static int count(List list, E o) {
    int counter = 0;
    for (int i=0; i<data.size(); i++) {
        E obj = data.get(i);
        if (obj.equals(o)) counter++;
     }
    return counter;
}</pre>
```

 Answer: No, this method is very inefficient for Linked Lists: get(i) always starts from the head and this is an O(n<sup>2</sup>) loop

# Efficient solutions are fundamentally different for:

y List	• Linked List
ount (E element){	<pre>int count (E element){</pre>
counter = 0;	int counter = 0;
(int i=0; i <size; i++){<="" td=""><td>Node finger = head;</td></size;>	Node finger = head;
<pre>if(data[i].equals(element)       counter ++;</pre>	<pre>while(finger != null){     if(finger.data.equals(element)</pre>
urn counter;	<pre>counter ++; finger = finger.next; Using while loop and next return counter; }</pre>
	<pre>y List ount (E element){ counter = 0; (int i=0; i<size; i++){<br="">if(data[i].equals(element) counter ++; urn counter;</size;></pre>

- But the principle of ADT forbids the use of underlying data structures directly!
- We need a uniform interface to iterate over List elements efficiently

# Efficient uniform iteration over List

- Problem: Efficient and uniform dispensing of values from the underlying data structures
- Solution: We create and use the common interface for iteration

# Extending operations for List ADT

- get()
- indexOf()
- add()
- remove()
- size()
- isEmpty()
- clear()
- contains()

But also method for efficient data traversal

> iterator()

## Iterator interface

- Iterators provide support for efficiently visiting all elements of an underlying data structure
- We customize the implementation of the iterator depending on the data structure
- We abstract away the details of how to access elements

#### public interface Iterator<E> :

boolean hasNext() - are there more elements for iteration?

E *next()* – return next element

# Example: Iterator for Array List

```
Can be a part of the ArrayList class
```

```
private class ArrayListIterator implements Iterator{
    ArrayList list;
    int nextIndex;
    public ArrayListIterator (ArrayList list){
          this.list = list;
          this.nextIndex = 0;
     }
    public boolean hasNext (){
          return (this.nextIndex < list.size());</pre>
    }
    public Object next(){
```

```
return list.data[nextIndex++];
```

```
}
```

}

```
private class ArrayListIterator implements Iterator{
    ArrayList list; Reference to the
    int nextIndex; actual Array List
     public ArrayListIterator (ArrayList list){
          this.list = list;
                                     We set it in the
          this.nextIndex = 0;
                                     constructor
     }
     public boolean hasNext (){
          return (this.nextIndex < list.size());</pre>
     }
    public Object next(){
          return list.data[nextIndex++];
     }
```

```
private class ArrayListIterator implements Iterator{
    ArrayList list;
                         Stores the current state of the iteration: the
     int nextIndex; position in the array to be returned next
     public ArrayListIterator (ArrayList list){
           this.list = list;
           this.nextIndex = 0;
     }
     public boolean hasNext (){
           return (this.nextIndex < list.size());</pre>
     }
     public Object next(){
           return list.data[nextIndex ++];
     }
```

```
private class ArrayListIterator implements Iterator{
    ArrayList list;
     int nextIndex;
     public ArrayListIterator (ArrayList list){
           this.list = list;
          this.nextIndex = 0;
                                          As long as
     }
                                          nextIndex is within
                                          valid bounds
     public boolean hasNext (){
           return (this.nextIndex < list.size());</pre>
     }
     public Object next(){
          return list.data[nextIndex++];
     }
}
```

```
private class ArrayListIterator implements Iterator{
     ArrayList list;
     int nextIndex;
     public ArrayListIterator (ArrayList list){
           this.list = list;
           this.nextIndex = 0;
     }
     public boolean hasNext (){
           return (this.nextIndex < list.size());</pre>
     }
                                          Return the element at position
                                          nextIndex, and advance
     public Object next(){
                                          nextIndex to the next position
           return list.data[nextIndex++];
     }
```

ArrayList *iterator*() returns array-specific Iterator:

public class ArrayList {
 Object[] data;
 int size;

public Iterator iterator (){
 return new ArrayListIterator(this);
}

## Iterator for Linked List

private class LinkedListIterator implements Iterator{

```
LinkedList list;
```

Node next;

 Same as before: reference to the actual Linked List

```
public LinkedListIterator (LinkedList list){
    this.list = list;
    this.next = list.head;
}
```

```
public boolean hasNext (){
}
```

```
public Object next(){
}
```

}

## Iterator for Linked List

private class LinkedListIterator implements Iterator{

```
LinkedList list;
```

Node next;

Stores the current state of the iteration: node to be read next

```
public LinkedListIterator (LinkedList list){
    this.list = list;
    next = list.head;
}
```

```
public boolean hasNext (){
}
```

```
public Object next(){
}
```

}

# Linked List Iterator: hasNext()

Which of the following is the correct implementation of hasNext()?

```
A. boolean hasNext(){
    return (this.list.size()>0)
}
```

```
B. boolean hasNext(){
        return (next.next != null)
    }
```

```
C. boolean hasNext(){
        return (next!= null)
    }
```

D. None of the above



# Linked List Iterator: next()

Which of the following is the correct implementation of *next()*?

```
A. Object next(){
        return this.list.get(next)
    }
```

```
B. Object next(){
        next = next.next;
        return next.data;
    }
```

```
C. Object next(){
        Object result = next.data;
        next = next.next;
        return result;
    }
```



D. None of the above

## Iterator for Linked List

```
public class LinkedListIterator implements Iterator{
   LinkedList list;
   Node next;
   ...
   public boolean hasNext (){
        return (next != null);
   }
}
```

```
public Object next(){
    Object result = next.data;
    next = next.next;
    return next;
}
```

}

### Linked List with its own iterator

public class LinkedList {

Node head; int size;

}

public Iterator iterator (){
 return new LinkedListIterator(this);
}

Uniform Counting with iterator() Works for both Array List and Linked List

public int count (List list, Object o) {
 int counter = 0;
 Iterator iter = list.iterator();
 while (iter.hasNext())
 if(o.equals(iter.next())) counter++;
 return counter;

}

#### Iterators: notes

- Iterator objects provide a common interface for traversing List ADT
- They have access to internal data representations
- They also store the state of traversal
- To implement an efficient iterator you need to **understand the mechanics of the underlying data structure**