

# Data structures: motivation

- ❑ The choice of a suitable data structure can make all the difference between an efficient and a failing program
- ❑ The input and output of any algorithm is stored inside a data structure
- ❑ Data structures organize data for quick and efficient access

# Examples of data structures

- ❑ Simple: *lists*, *stacks* and *queues*
- ❑ More intricate - but still very useful: *heaps*, *search trees*, *hash tables*, *Bloom filters*, *union-find* ...

- ❑ Why so many?

Because different data structures support different sets of operations and are good for different types of tasks

# Know what exists and what it is good for

- ❑ We will discuss the pros and cons of each data structure for a particular task
- ❑ The fewer operations the data structure supports - the faster the operations will be

Think about the operations that you need a data structure to support



Choose the best data structure - the one that supports only required operations, and not more.

# 4 levels of Data Structure Proficiency

- ❑ Level 0: **ignorance**
- ❑ Level 1: **cocktail party awareness**
- ❑ Level 2: **solid literacy**: know which data structures are appropriate for which types of tasks and comfortable using them
- ❑ Level 3: **hardcore** programmers and computer scientists: understand the internals of existing and implement new data structures



# **Basic Data Structures: Arrays and Linked Lists**

Lecture 02.01 by Marina Barsky

# ARRAYS

```
long arr[] = new long[5];
```

Java

```
long arr[5];
```

C/C++

```
arr = [None] * 5
```

Python

1	5	17	3	25
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1D

1	5	17	3	25
8	2	36	5	3

2D

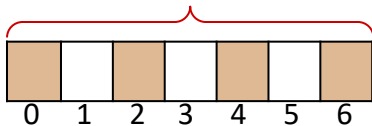
# Definition

Array:

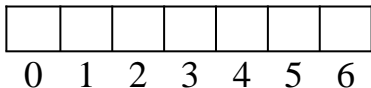
**Contiguous** area of memory containing **equal-size** elements indexed by contiguous integers.

The maximum number of elements that can fit into the allocated memory is called a *capacity* of the array.

The number of elements currently in the array is called a *size* of the array.



# What's Special About Arrays?



**Constant-time access** to any element by index  $i$ .

Computed as:

$$\text{array\_addr} + \text{elem\_size} \times (i)$$



# Multi-Dimensional Arrays

```
int arr [3][6];
```

(0,0)					
			(2,3)		

$\text{array\_addr} + \text{elem\_size} \times (2 \times 6 + 3)$

(0, 0)
(0, 1)
(0, 2)
(0, 3)
(0, 4)
(0, 5)
(1, 0)
.

# Arrays: Time for Common Operations

	Add	Remove
Beginning		
End		
Middle		

5	8	3	12			
---	---	---	----	--	--	--

size=4

# Arrays: Time for Common Operations

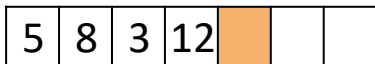
	Add	Remove
Beginning		
End	$O(1)$	
Middle		

5	8	3	12	4		
---	---	---	----	---	--	--

$A[4] = 4$   
size = 5

# Arrays: Time for Common Operations

	Add	Remove
Beginning		
End	$O(1)$	$O(1)$
Middle		



Remove  $A[4]$   
size=4

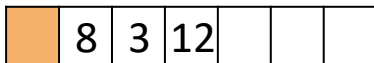
# Arrays: Time for Common Operations

	Add	Remove
Beginning		
End	$O(1)$	$O(1)$
Middle		

5	8	3	12			
---	---	---	----	--	--	--

# Arrays: Time for Common Operations

	Add	Remove
Beginning		
End	$O(1)$	$O(1)$
Middle		



# Arrays: Time for Common Operations

	Add	Remove
Beginning		
End	$O(1)$	$O(1)$
Middle		

8		3	12			
---	--	---	----	--	--	--

# Arrays: Time for Common Operations

	Add	Remove
Beginning		
End	$O(1)$	$O(1)$
Middle		

8	3		12			
---	---	--	----	--	--	--



# Arrays: Time for Common Operations

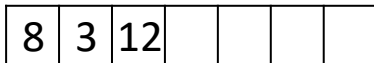
	Add	Remove
Beginning		$O(n)$
End	$O(1)$	$O(1)$
Middle		

8	3	12				
---	---	----	--	--	--	--

size=3

# Arrays: Time for Common Operations

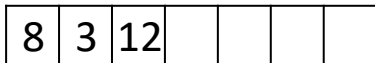
	Add	Remove
Beginning	$O(n)$	$O(n)$
End	$O(1)$	$O(1)$
Middle		



size=3

# Arrays: Time for Common Operations

	Add	Remove
Beginning	$O(n)$	$O(n)$
End	$O(1)$	$O(1)$
Middle	$O(n)$	$O(n)$



size=3

# Summary: Arrays

- ❑ Array: **contiguous area** of memory consisting of **equal-size elements** indexed by contiguous integers
- ❑ **Constant-time access** to any element **by** location (**index**)
- ❑ **Constant time** to add/remove **at the end**
- ❑ **Linear time** to add/remove at an arbitrary location