

Greedy scheduling

Lecture 05.02

SAMPLE PROBLEM 4

Movie star scheduling

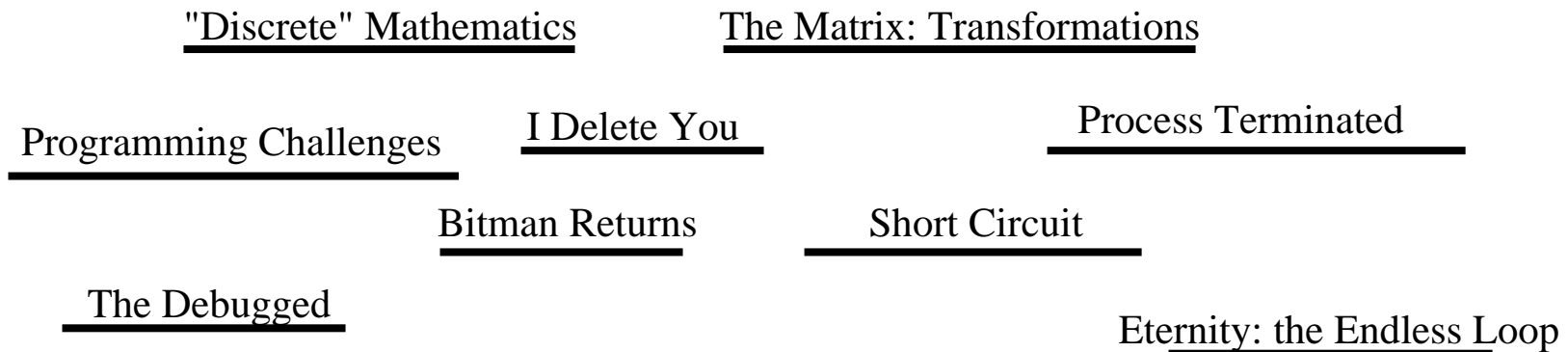
Movie Star Scheduling Problem

A movie star has been offered the leading role in several upcoming movies

They want to select the **maximum number of roles** such that no movies overlap in time

We call the movies that overlap *conflicting movies*

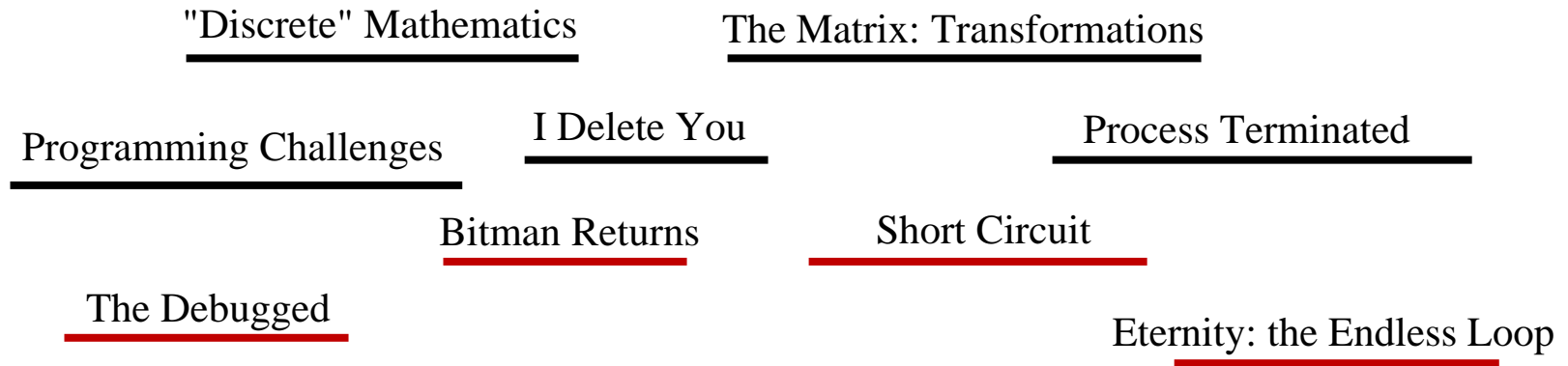
Movies with their start and end times shown as an interval of time:



What is an optimal solution for this problem instance?

Movie Star Scheduling Problem

Movies with their start and end times shown as an interval of time:



Optimal solution: 4 jobs

Movie Star Scheduling Problem

- This problem is known as *maximum independent set in an interval graph*.
- Each interval has a start and end value.

Movie Scheduling Problem

Input: A set I of n intervals on the line.
Output: The largest subset of non-conflicting intervals which can be selected from I

How long will exhaustive computation take?

Movie Star Scheduling Problem

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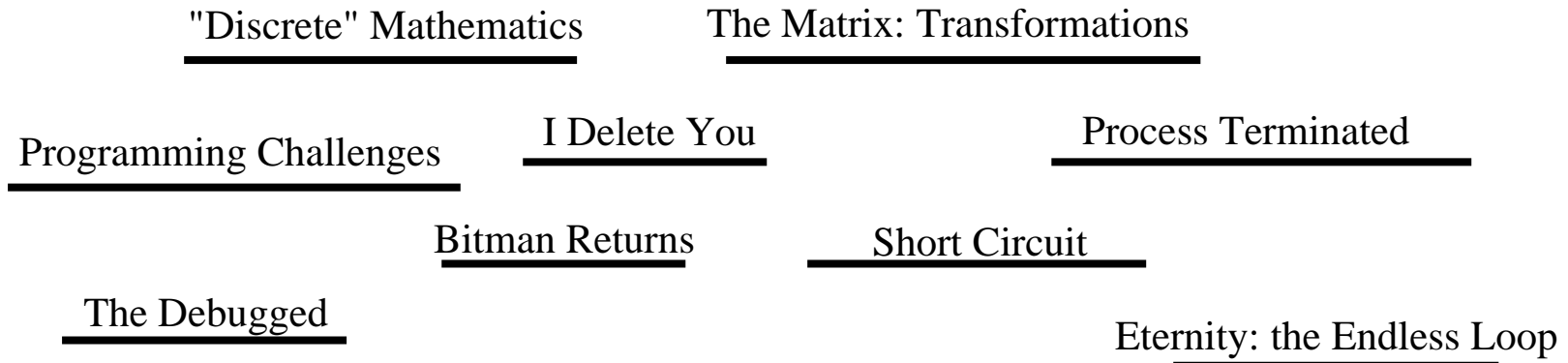
Movie Scheduling Problem

Input: A set I of n intervals on the line.
Output: The largest subset of non-conflicting intervals which can be selected from I

2ⁿ different interval subsets: similar to knapsack 01 – the interval is either selected or not

Greedy move: version 1

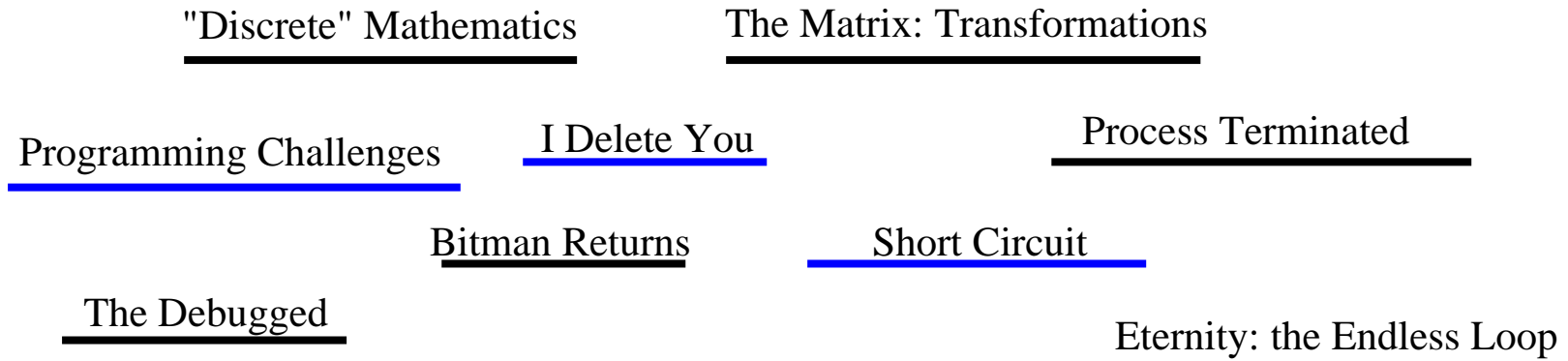
1. *Starting-First*. Accept the job that starts soonest and doesn't conflict.



What jobs will be selected by the starting-first algorithm?

Greedy move: version 1

1. *Starting-First*. Accept the job that starts soonest and doesn't conflict.



What jobs will be selected by the starting-first algorithm?

Is "starting-first" a safe move?

Can we miss an optimal solution?

Counterexample for Starting-First



Starting-first is not a safe move ■

Greedy move: version 2

2. *Shortest-First*. Accept the shortest job that doesn't conflict.

"Discrete" Mathematics

The Matrix: Transformations

Programming Challenges

I Delete You

Process Terminated

Bitman Returns

Short Circuit

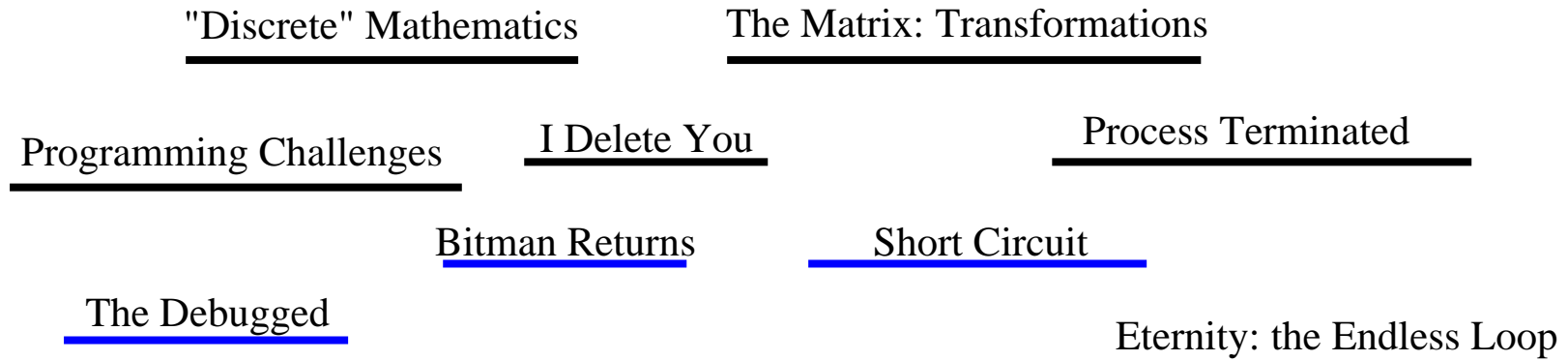
The Debugged

Eternity: the Endless Loop

What jobs will be selected by the shortest-first algorithm?

Greedy move: version 2

2. *Shortest-First*. Accept the shortest job that doesn't conflict.



What jobs will be selected by the shortest-first algorithm?

Is "shortest-first" a safe move?

Can we miss an optimal solution?

Counterexample for shortest-first



Shortest-first is not a safe move ■

Movie Star Problem

Input: A set I of n intervals on the line.
Output: The largest subset of conflicting intervals which can be selected from I

Maybe greedy approach does not work here?

Maybe we need to do an exhaustive search?
Or use another strategy?

Greedy move: version 3

3. *Ending-First*. Accept the job that ends soonest and doesn't conflict

"Discrete" Mathematics

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What jobs will be selected by the ending-first algorithm?

Greedy move: version 3

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What jobs will be selected by the ending-first algorithm?

Is "ending-first" a safe move?

Can we miss an optimal solution?

Theorem

Ending-first is a safe move

Proof

- Let x be a job which contains the end point which is leftmost among all remaining intervals.
- Other jobs may well have started before x , but all of these must conflict with x , so we can select at most one from the group.
- The first of these jobs to terminate is x , so any of the other conflicting jobs potentially block out more opportunities to the right of it.
- Clearly we can never lose by picking x . ■

The Matrix: Transformations



Process Terminated



x



Leaves most non-conflicting jobs to choose from

The Matrix: Transformations

Process Terminated

x

Leaves most non-conflicting
jobs to choose from

- Note that the proof of this theorem did not use an exchange argument
- Instead, we use a *lower-bound argument*: we argue that any solution will be no better without this greedy choice
- Our greedy choice achieves the lower bound of all possible choices

Read more about Greedy Algorithms

10.1 and 10.2 of the textbook

Puzzle: Bridge Crossing at Night

- A group of 4 people with 1 flashlight need to cross a rickety bridge at night.
- A maximum of 2 people can cross the bridge at one time, and any party that crosses (either 1 or 2 people) must have the flashlight with them.
- A pair must walk together at the rate of the slower person's pace.
- The flashlight must be walked back and forth - it cannot be thrown.

The time taken by each person:

Ann takes 1 minute

Bob takes 2 minutes

Cat takes 5 minutes

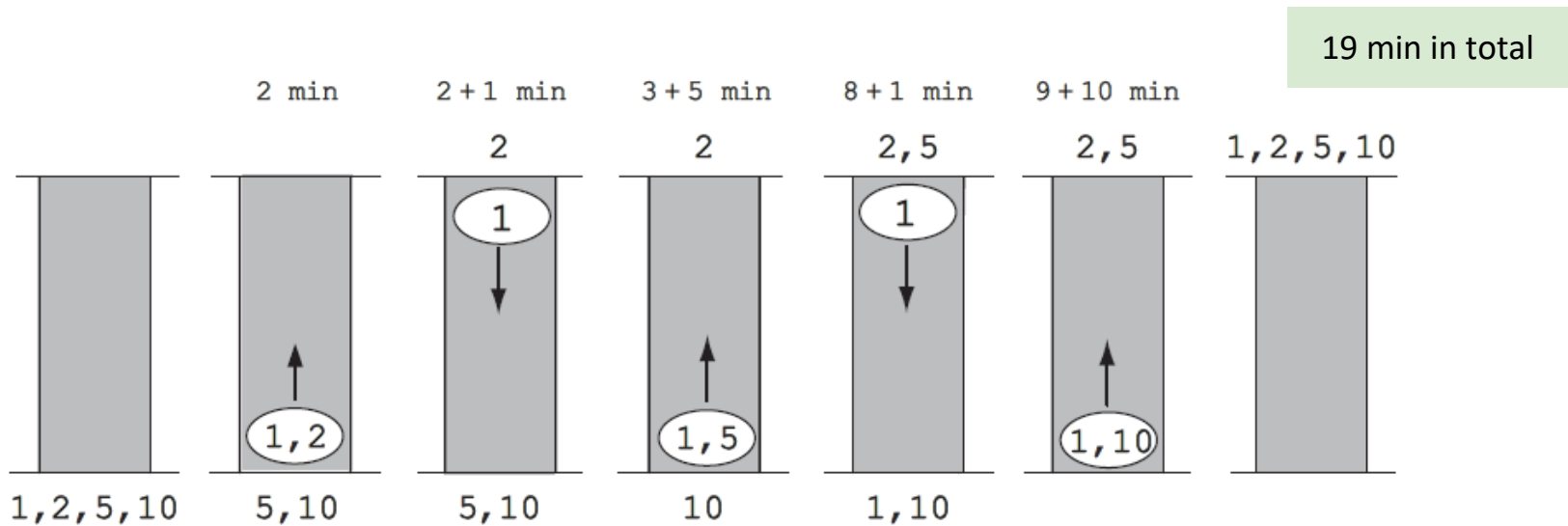
Don takes 10 minutes

Find the fastest way they all can cross the bridge

Greedy Algorithm

- 1. Ann takes 1 minute
- 2. Bob takes 2 minutes
- 5. Cat takes 5 minutes
- 10. Don takes 10 minutes

Always send 2 fastest people available, and always send the fastest person back to return the light



Did the algorithm find an optimal solution?

How to reduce the total time?

Play here:

<https://www.inwebson.com/demo/cross-the-bridge/>

What is the main idea of the optimal solution?
What is a greedy move?